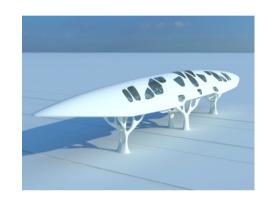
Outer-Point Generation – A Decomposition Algorithm for solving MINLPs



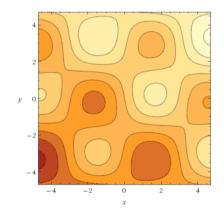
Ivo Nowak

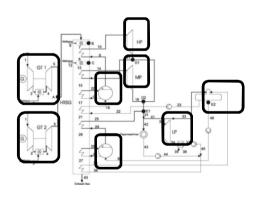


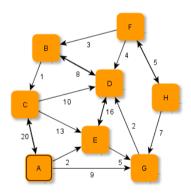












SCIP Workshop, Berlin, 1.10.2014

Main objectives of this talk

Present a new decomposition method, called Outer-Point
 Generation combining MINLP and Column Generation

2. Discuss ideas how this method can be implemented in SCIP

Agenda

Motivation: Path Based Column Generation

MINLP Outer-Point Generation

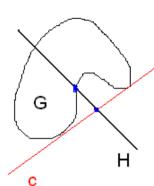
Implementation in SCIP

Next Steps

Classical decomposition methods are efficient, if the duality gap is small

Quasi-separable reformulation and convex relaxation

optimization problem (P): $c^* := \min\{c(x) : x \in G \cap H\}$ convex relaxation (C): $c_{\mathsf{relax}} := \min\{c(x) : x \in \mathsf{conv}(G) \cap H\}$ duality gap: $\mathsf{gap} := c^* - c_{\mathsf{relax}}$



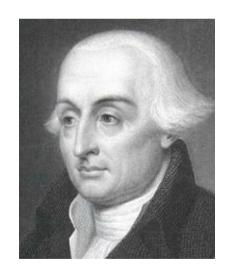
duality gap

Four possibilities for solving the convex relaxation (C)

- 1) Lagrangian Decomposition (extreme points of G)
- 2) Column Generation (Dantzig-Wolfe Decomp, inner points of G)
- 3) Outer Approximation (cutting planes of G)
- 4) Frank-Wolfe Algorithm (Linear approximation of the objective)

If gap is small, then solution of (C) is a good starting point for (P)

Planning & Control: gap typically small gap typically large



Joseph-Louis Lagrange 1736 -1813

xOPT is a parallel path-based column generation framework for crew and aircraft scheduling developed at Lufthansa Systems

Crew rostering problem:

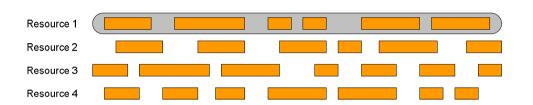
variables: 772.927.392

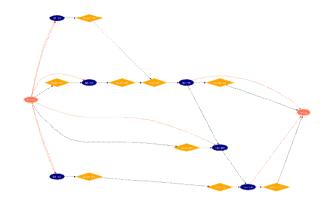
(pricing) sub-problems: 3.012

■ coupling constraints: ~500.000

This problem cannot be solved with a generic optimization tool in reasonable time

The xOPT column generation approach for weakly nonlinear network problems





Constraint Shortest Path pricing problems

LP/IP master problem

 $\label{eq:minimum} \begin{array}{lll} \mbox{Min} & L_k(p,\pi) & \mbox{CSP} \\ & h_k(p) \leq 0 & \mbox{(nonlinear constraints)} \\ & p \in P_k & \mbox{(network constraints)} \end{array}$

Min
$$c(x,s)$$
 LP/IP
$$Ax = b+s \text{ (linear constraints)}$$

$$x \in X, s \in S$$

CSP-Solver:

- Label-based dynamic programming algorithm with relaxed dominance
- Dynamic network filtering and expansion

LP-Solver

- Inexact bundle method
- Interior point start heuristic

IP-Solver:

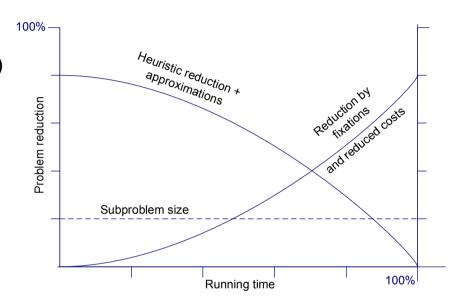
- Column/constraint fixing by rapid branching based on
- perturbation heuristic

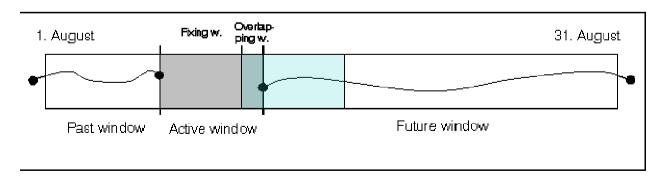
Chart 6 Outer-Point Generation

The dynamic Reduce and Generate approach keeps the size of sub-problems manageable

- Search-space reduction by a (heuristic) active set strategy regarding columns, constraints, nodes and arcs
- Few main iterations, e.g. 20

Step-wise fixing by moving a timewindow over the optimization period





Agenda

Motivation: Path Based Column Generation

MINLP Outer-Point Generation

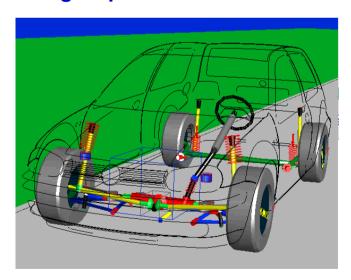
Implementation in SCIP

Next Steps

Goal: Build a new point-based decomposition method combining MINLP and CG

- The new method shall be used for **industry optimization problems**, which are
 - sparse (-> decomposition)
 - strongly nonlinear, nonconvex, many variables
 - **simulation-based**, i.e. no algebraic formulation available
- Currently solved by heuristics (e.g. evolutionary algorithms)

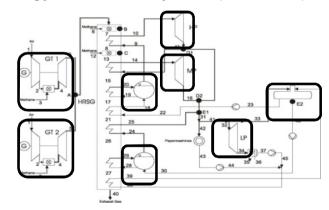
Design Optimization



Quelle: http://www.simpack.com

Planning

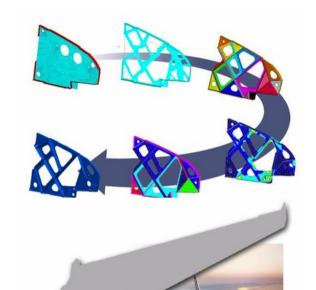
Energy conversion system (TU-Berlin)



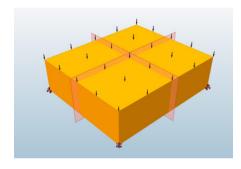
Multidisciplinary Optimization (MDO)

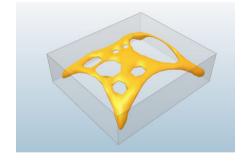
Outer-Point Generation

Decomposition of structural optimization problems



Airbus A380, structural optimization of a wing's rib

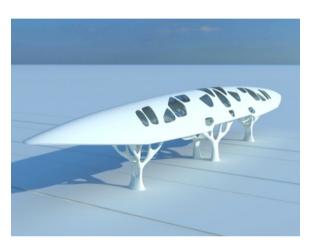




Design Space subjected to topology optimization







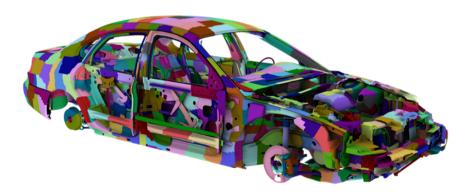


Pegasus bridge, final re-design

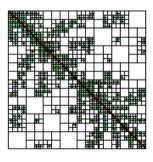
Decomposition of sparse optimization problems

Quasi-separable reformulation by

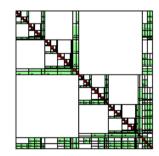
- partition problem into sub-problems (domain decomposition)
- connect sub-problems by coupling constraints: Ax = 0
- (P) $\min c(x), \qquad Ax = 0, \quad (x_k, y_k) \in G_k, \quad k \in K$ $G_k := \{x \in X_k : g_k(x) \le 0\}$ $X_k \quad \text{is a mixed integer interval}$



http://industry.it4i.cz/en/research/basic-research/feti-based-domain-decomposition-methods/



sparse-structure

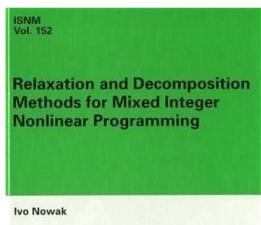


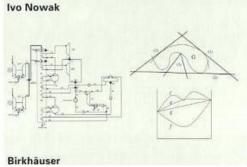
block-structure

(e.g. of a FEM-system)

Preliminary experiments with the MINLP-solver LaGO 2004

- Preliminary implementation of Lagrangian Decomposition (LD) for MINLP within LaGO together with Stefan Vigerske
- Experiments of LD with QQPs and MINLPlib showed:
- LD was not competitive to branch-and-cut because:
 - MINLPs have typically a large duality gap.
 - No special method for efficiently solving MINLP sub-problems was used.





What are the key-properties of the CG approach?

How should they be adapted for solving MINLPs?

1. No huge branch-bound trees

CG: rapid branching (works if the duality gap is small)

MINLP: Claim 1: Use a nonconvex master problem

2. Efficient sub-problem solving

CG: (1) highly tuned algorithm CSP/DP for enumerating all reduced cost solutions

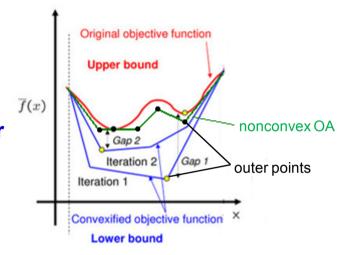
(2) dynamic search space reduction

MINLP : Claim 2: Fast reduction of the approximation error

The new Outer-Point Generation (OG) approach

Claim 1: Use a nonconvex master problem

- Idea:
 - Define the master problem by a nonconvex point-based Outer Approximation (OA) based on Piecewise Linear Envelopes
 - Improve the OA by generating tighter outer-points (vertices) of the OA

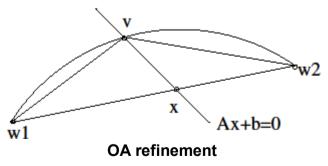


v: projected solution

x: OA-solution

Claim 2: Fast reduction of the approximation error

- Idea:
 - Use a two-phase approach for reducing the OA-error:
 - (1) Column generation
 - (2) Sequential approximation: (OA-refinement)
 projection of ε-OA-solutions onto the feasible set



conv $\{w1,w2\} \rightarrow conv \{w1,v\} + conv \{v,w2\}$

Remarks

Why OG and not Branch-Cut-And-Price?

Cut-based OA is improved by branching (which we want to avoid)

(for integer variables)

■ Point-based OA can be improved without branching (for 'nonlinear' variables)

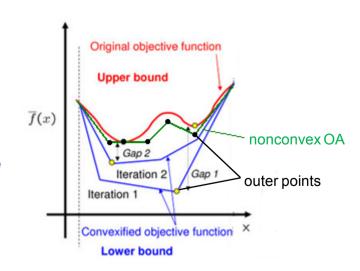


Active set strategy for reducing points (and cuts) is neccessary



Derivative-Free Black-Box Optimization

- Sequential approximation of surrogate models of black-box functions
- Refine surrogate models at points generated by the OG-algorithm



Initialization: Outer-Approximation by Piecewice Linear Envelopes

Piecewise linear Outer-Approximation:

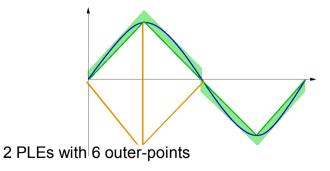
(OA)
$$\min c(x), \qquad Ax = 0, \quad (x_k, y_k) \in \hat{G}_k, \quad k \in K$$
 $\hat{G}_k := \bigcup_{j \in J_k} \operatorname{conv}(V_{kj})$ x: primary (coupling) variables y: secondary (nonlinear) variables

y: secondary (nonlinear) variables

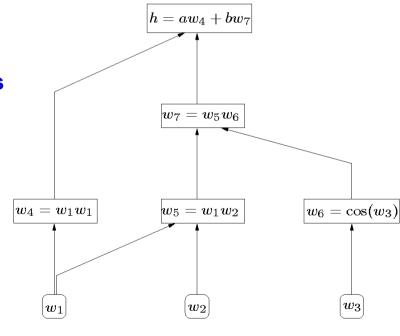
Branch-and-Refine approach of S. Leyffer, A. Sartenaer, and E. Wanufelle, 2008:

Decompose nonlinear functions into bilinear and univariate expressions using expression trees

Replace bilinear and univariate expressions by PLEs (Piecewise Linear Envelopes) using outer-points







$$g(x)=ax_1^2 + bx_1x_2\cos(x_3)$$

Refinement of Piecewise Linear Envelopes (PLEs)

at an interpolation point function The function function

- Outer-points of a PLE are chosen such that the approximation error is minimized
- OA-error depends on size of linear pieces of the PLEs
- **Continuity of g**: If the OA-error is small enough, then the **solution of (P)** is in the **neighborhood** of an ϵ -OA-solution (-> ϵ -solving the OA gives the solution)

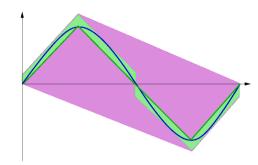
Phase 1: Column Generation

Goal: Reduce the OA-error up to the duality gap (optional)

LP - master problem in the space of primary variables x

(LP-MP)
$$\min c(x), \qquad Ax = 0, \quad x_k \in \text{conv}(W_k), \quad k \in K$$

$$W_k := \{\hat{x}_{ki} : i \in I_k\} \quad (\hat{x}_{ki}, \hat{y}_{ki}) \in V_k$$



Column generation

- 1. Compute a **dual solution** of (LP-MP)
- 2. Compute ε-solutions (columns) of the MINLP pricing sub-problems regarding a reduced cost direction
- 3. Local refinement of PLEs in the neighborhood of sub-problem solutions

Phase 2: Sequential Approximation

Goal: Reduce the OA-error up to ε by improving Lower and Upper Bounds

(MIP-MP)
$$\min c(x), \qquad Ax = 0, \quad x_k \in \bigcup_{j \in J_k} \operatorname{conv}(W_{kj}), \quad k \in K$$

$$W_{kj} := \{\hat{x}_{ki} : i \in I_{kj}\} \quad (\hat{x}_{ki}, \hat{y}_{ki}) \in V_k$$

Sequential Approximation (Projection)

- 1. Enumerate ε-solutions of MIP-MP (OA) by Branch & Cut (and save valid cuts for warm start)
- For every ε-OA-solution: (LB-improvement)
 - a) project the OA-solution onto the feasible set G
 - b) local refinement of PLEs in the neighborhood of the projected solution
- 3. Compute a (or several) **local minimizer(s)** of problem (P) starting from the ε-OA-solution(s) (UB-improvement)
- 4. Local refinement of PLEs in the neighborhood of the minimizers

Solving the master problems

MIP master problem solver for Step 1:

(in the primary space)

- Branch&Cut
- Many outer points -> active-set strategy:
 - use a **bundle of points** of the OA which are **active** in LPs of **B&B partitions**
 - use global cuts which are computed by aggregating B&B cuts

MINLP local search for Step 3:

(in the original space)

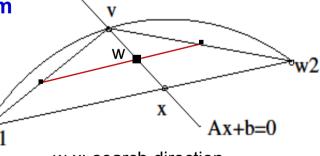
NLP or feasibility pump

or search direction by solving a perturbed problem defined by shifting active outer-points towards the projected OA solution

UB-LB gives ε

w: solution of perturbed OA

v: projected OA solution

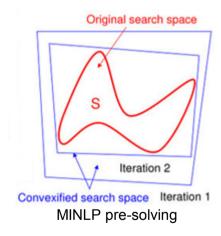


w-x: search direction

Solving the sub-problems

Two MINLP sub-problem solvers for pricing (CG) and projection (Step 3) (in the primary and secondary space)

- Start with MINLP pre-solving where cuts are stored for warm start
- NLP or feasibility pump for computing UB
 - starting from ε-solutions of sub-OA
- MINLP-Branch&Cut and/or Sequential Approximation



CG versus OG

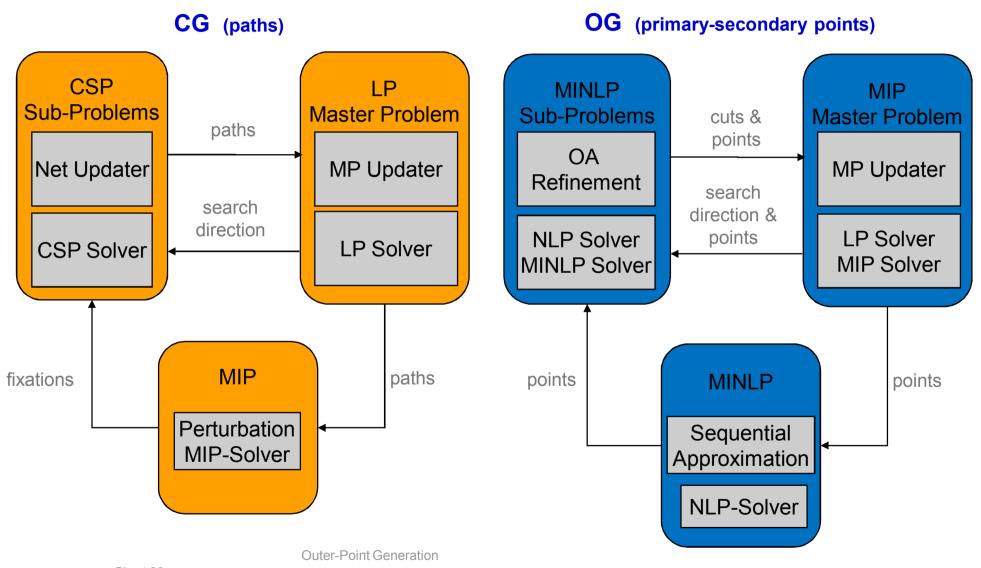


Chart 22

Agenda

Motivation: Path Based Column Generation

MINLP Outer-Point Generation

Implementation in SCIP

Next Steps

Proposal for the development of a decomposition framework

- 1. CG for MIQQPs using SCIP/GCG and SCIP/MINLP as a pricing-solver
 - unit test, local search, e.g. using feasibility pump or new perturbation heuristic
- 2. Parallel decomposition framework using a MIP master problem
 - for MIQQP
 - based on GCG

3. Extension to MINLPs

- automatic quasi-separable reformulation by partitioning the sparsity graph
- point-based OA using PLEs

Agenda

Motivation: Path Based Column Generation

MINLP Outer-Point Generation

Implementation in SCIP

Next Steps

Next Steps

- We presented the OG algorithm with similar properties as CG
- Discussion on algorithmic details, e.g. active set strategy, master and sub-solvers,..
- Finishing internal report "The Outer-Point Generation Algorithm"
- Implementation of a prototype for a specific application
 - e.g. a bilinear problem

Applications

- engineering optimization (HAW)
- transport (ZIB)
- energy conversion (TU-Berlin)

Literature

- I. Nowak, *Parallel Decomposition Methods for Nonconvex Optimization Recent Advances and New Directions*, MAGO, Malaga, Spain, 2014
- R. Borndörfer, A. Löbel, M. Reuther, T. Schlechte, and S. Weider. *Rapid branching*. Public Transport, 2013
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- G. Gamrath and M. Lübbecke. *Experiments with a generic Dantzig-Wolfe decomposition for integer programs*. In Festa 2010.
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- I. Nowak, Relaxation and Decomposition Methods for Mixed Integer Nonlinear Programming, Birkhäuser Basel, 2005

Thank you!